

# Stochastic Knapsack Problem with random weights

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## The Deterministic Knapsack Problem

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## The Stochastic Knapsack Problem with Random Weights

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- $\chi_i$ : **independently normally distributed** weight of item  $i$

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- $\mu_i, \sigma_i$ : mean and standard deviation of  $\chi_i$



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- $\mu_i, \sigma_i$ : mean and standard deviation of  $\chi_i$
- $x \in \{0, 1\}^n$ : decision vector

# Outline

- 1 Introduction
- 2 Stochastic Knapsack Problem with simple recourse
  - Problem Formulation
  - Analytic description
- 3 Problem Solving Method
  - Relaxed Stochastic Knapsack Problem
  - Approximation by convolution
  - Branch-and-Bound Algorithm

## Stochastic Knapsack Problem with simple recourse (SRKP)

$$\max_{x \in \{0,1\}^n} \mathbb{E} \left[ \sum_{i=1}^n r_i \chi_i x_i \right] - d \cdot \mathbb{E} \left[ \left[ \sum_{i=1}^n \chi_i x_i - c \right]^+ \right]$$

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- $d > 0$ : penalty factor per weight unit



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$$J_{det}(x) =$$

$$\sum_j r_j \mu_j x_j - d \cdot \left[ \hat{\sigma} \cdot f\left(\frac{c - \hat{\mu}}{\hat{\sigma}}\right) - (c - \hat{\mu}) \cdot \left[ 1 - F\left(\frac{c - \hat{\mu}}{\hat{\sigma}}\right) \right] \right]$$

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- Branch-and-Bound algorithm
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- Apply "Approximation by convolution" - method to approximate the gradient of the objective function

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## Definition

*Relaxed Stochastic Knapsack Problem (with simple recourse):*

$$\max_{x \in [0,1]^n} \mathbb{E} \left[ \sum_{i=1}^n r_i \chi_i x_i \right] - d \cdot \mathbb{E} \left[ \left[ \sum_{i=1}^n \chi_i x_i - c \right]^+ \right]$$

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Update  $x^k$  as follows:

$$x^{k+1} = x^k + \epsilon^k r^k$$

where  $r^k = \nabla_x j(x, \chi)$  and  $(\epsilon^k)_{k \in \mathbb{N}}$  is a  $\sigma$ -sequence

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For all  $i = 1, \dots, n$ :

If  $x_i^{k+1} > 1$  set  $x_i^{k+1} = 1$  / If  $x_i^{k+1} < 0$  set  $x_i^{k+1} = 0$

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Then, for small values of  $r > 0$ , we get the following approximation of a locally integrable real valued function  $f$ :

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$$h := \frac{3}{4}(1 - x^2)\mathbb{1}_1(x)$$

( $\mathbb{1}_1$ : indicator function of the interval  $] -1, 1[$ )

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## Approximated gradient

$$\nabla_x j(x, \chi) = \begin{pmatrix} r_1 \chi_1 \\ \cdot \\ \cdot \\ \cdot \\ r_n \chi_n \end{pmatrix} + d \cdot \left( \frac{3}{4r} \left( 1 - \left( \frac{g(x, \chi)}{r} \right)^2 \right) \mathbb{1}_1 \left( \frac{g(x, \chi)}{r} \right) \chi \cdot g(x, \chi) - \mathbb{1}_{\mathbb{R}^+}(g(x, \chi)) \cdot \chi \right)$$



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where  $g(x, \chi) := \sum_{i=1}^n \chi_i x_i$

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## Motivation

$$\frac{\partial J_{det}}{\partial \hat{\sigma}}(x) =$$

$$\hat{\sigma} := \sqrt{\sum_{i=1}^n \sigma_i^2 x_i^2}$$

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## Motivation

$$\frac{\partial J_{det}}{\partial \hat{\sigma}}(x) = -d \cdot f\left(\frac{c - \mu}{\sigma}\right) < 0$$

$$\hat{\sigma} := \sqrt{\sum_{i=1}^n \sigma_i^2 x_i^2}$$



## The Branch-and-Bound Algorithm

Ranking

## The Branch-and-Bound Algorithm

### Ranking

- Number of objects dominated
- value of  $\frac{r_i^2}{\sigma_i}$

## The Branch-and-Bound Algorithm

Ranking

Plunging

## The Branch-and-Bound Algorithm

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- Beginning at the root, add current item iff objective function increases
- $INF \leftarrow$  maximum value of the objective function found
- Add branch found to list of waiting branches;  
set assigned value  $SUP$  to  $\infty$

## The Branch-and-Bound Algorithm

Ranking

Plunging

Branch choosing

## The Branch-and-Bound Algorithm

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- **If** no branch left on list of branches  $\rightarrow$  step G.
- **Else** take branch having maximum objective function value.  $\rightarrow$  step D.

## The Branch-and-Bound Algorithm

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Branch choosing

Delete bad branches

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- **If**  $SUP > INF \rightarrow$  step E.
- **Else** delete branch.  $\rightarrow$  step C.



## The Branch-and-Bound Algorithm

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- **If** no accepted item left in selected branch that does not already have a plunged or rejected subtree, delete branch from list. → step C.
- **Else** choose first accepted item that does not already have a plunged or rejected subtree. Calculate upper bound SUP for the subtree defined by rejecting this item. → step F.

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- **If**  $SUP \leq INF$ , reject subtree,  $\rightarrow$  E.
- **Else** plunge subtree as described in B, add the found branch with value SUP to list. Update INF.  $\rightarrow$  step C.

## The Branch-and-Bound Algorithm

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The current value INF is the optimal solution.

# Thank you!

# Merci!

# Danke!

n	Approximation by convolution					Cohn/Barnhard: "The Stochastic Knapsack Problem with Random Weights"				
	Upper Bound	CPU-time (msec) continuous	Optimum	considered nodes	CPU-time (sec) B-and-B	Upper Bound	CPU-time (msec) continuous	Optimum	considered nodes	CPU-time (sec) B-and-B
C./B.	4676.208	4	4618.025	100	0.342	4759.000	0	4618.025	144	0.000
15	4934.583	4	4889.781	41	0.139	5146.927	0	4889.781	65	0.002
20	6690.744	6	6650.513	80	0.348	6936.017	0	6650.517	280	0.003
30	10279.908	9	10264.683	455	2.808	10529.541	0	10264.756	2525	0.037
50	16954.343	12	16950.579	13173	131.171	17224.803	0	16950.757	364960	779.325
75	25519.688	16	25513.555	63972	934.550	25811.775	0	*	*	*
100	33846.095	22	*	*	*	34131.754	0	*	*	*
150	50607.008	31	*	*	*	50932.104	0	*	*	*
250	85098.136	52	*	*	*	85459.649	1	*	*	*
500	170110.459	104	*	*	*	170503.708	3	*	*	*
1000	340922.966	240	*	*	*	340822.740	5	*	*	*
5000	1703811.095	1110	*	*	*	1704935.949	107	*	*	*

—Mean over 50 randomly created instances—

\* average CPU time exceeds 1h