Stochastic Optimization IDA PhD course 2011ht

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5. Lecture: Two-Stage Problems 10. November 2011





1 Randomness occurs in the constraint function



1 Randomness occurs in the constraint function

2 Simple Recourse problems - Useful Theorems



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3 Two-Stage Programming

- Example: TSKP
- Simple Recourse as a special case of Two-Stage problems
- Deterministic reformulation in case of discrete distributions
- Main theoretical result



Outline

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$\mathsf{Deterministic}\ \mathsf{Opt}.\ \mathsf{Model} \to \mathsf{Stochastic}\ \mathsf{Programming}\ \mathsf{Model}$

$$\begin{array}{ll} \max_{x \in X} \quad f(x) & \min_{x \in X} \quad f(x) \\ \text{s.t.} \quad G(x) \le 0 & _ & \text{s.t.} \quad G(x, \chi) \le 0 \end{array}$$

$$\begin{split} &\chi \in \Omega \subseteq \mathbb{R}^s: \text{ random vector} \\ &G: \mathbb{R}^n \times \mathbb{R}^s \to \mathbb{R}^m \\ &G(x,\chi) = (g_1(x,\chi), \dots, g_m(x,\chi)) \end{split}$$





- \blacksquare Feasible in all possible cases \rightarrow conservative/infeasible/trivial
- \blacksquare Feasible with high probability \rightarrow chance-constraint \checkmark
- \blacksquare Average violation not too bad \rightarrow simple recourse \checkmark



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- Feasible after correction has been made



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- \blacksquare Feasible with high probability \rightarrow chance-constraint \checkmark
- Average violation not too bad \rightarrow simple recourse \checkmark
- Feasible after correction has been made \rightarrow multi-stage



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Basic Idea(s)

- Allow violation
- Restrict average "amount of violation"
- Introduce penalty per "violation unit"
- Find good trade-off: cost \leftrightarrow penalty



Simple Recourse problems - Useful Theorems

Simple-Recourse problem



 $\min_{x\in X} f(x)$



$$\min_{x \in X} \quad f(x) + \sum_{i=1}^{m} d_i \cdot \mathbb{E}\left[[g_i(x, \chi)]^+ \right]$$

$$\bullet d_i > 0 \ \forall i \in \{1, \ldots, m\}$$



$$\min_{x\in X} f(x) + \sum_{i=1}^{m} d_i \cdot \mathbb{E}\left[\left[g_i(x,\chi)\right]^+\right]$$

■
$$d_i > 0 \ \forall i \in \{1, ..., m\}$$

■ $[x]^+ = max(0, x)$



Theorem Let • $\gamma(x) := \mathbb{E}\left[\left[g(x, \chi)\right]^+\right]$



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27	
• $\gamma(x) := \mathbb{E}\left[[g(x, \chi)]^+ \right]$	
• $\mathbb{E}[g(x,\chi)] < \infty \ \forall x \in \mathbb{R}^n$	



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Theorem Let • $\gamma(x) := \mathbb{E}[[g(x, \chi)]^+]$ • $\mathbb{E}[g(x, \chi)] < \infty \ \forall x \in \mathbb{R}^n$ Then the following holds: $g(\cdot, \hat{\chi})$ is convex $\forall \hat{\chi} \implies \gamma$ is convex



$$\min_{x\in X} f(x) + \sum_{i=1}^{m} d_i \cdot \mathbb{E}\left[\left[g_i(x,\chi)\right]^+\right]$$

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$$d_i > 0 \ \forall i \in \{1, ..., m\}$$

■ $[x]^+ = max(0, x)$



Let

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$$\gamma(x) := \mathbb{E}\left[\left[a(\chi)x - b\right]^+\right]$$



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- **1** γ is convex
- **2** For all $x \in \mathbb{R}^n \gamma(x) < \infty$.



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$$\gamma(x) := \mathbb{E}\left[\left[a(\chi)x - b\right]^+\right]$$

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Then the following holds:

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- **2** For all $x \in \mathbb{R}^n \gamma(x) < \infty$.
- $\mathbf{3} \gamma$ is Lipschitz continuous.

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- **4** γ is differentiable wherever Φ is continuous.

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Then the following holds:

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- **2** For all $x \in \mathbb{R}^n \gamma(x) < \infty$.
- $\mathbf{3} \gamma$ is Lipschitz continuous.
- **4** γ is differentiable wherever Φ is continuous.
- **5** $a(\chi)x b$ discretely distributed $\Rightarrow \gamma$ is piecewise linear.

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Feasibility in all scenarios: Conservative solutions



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- Chance-Constraints: Infeasibility of solution possible



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- Simple Recourse: No actual "recourse" action


Why Two-Stage Programming?

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- In the end: feasible solution



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- In the end: feasible solution \rightarrow most realistic?



Why Two-Stage Programming?

- Feasibility in all scenarios: Conservative solutions
- Chance-Constraints: Infeasibility of solution possible
- Simple Recourse: No actual "recourse" action
- \blacksquare In the end: feasible solution \rightarrow most realistic?
- Costs/rewards of recourse action taken into account





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- **2** (Pre)-decision taken based on (statistical) information of uncertain parameters



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- 3 Objective: Minimize expected total cost



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- (Pre)-decision taken based on (statistical) information of uncertain parameters
- 3 Objective: Minimize expected total cost
- 4 Between first and second stage: all random parameters come to be known
- **5** Second stage: deterministic problem
- 6 Make corrective decision ("recourse"):
 - \rightarrow Make solution feasible
 - ightarrow Decrease total cost





$$\min_{x\in X} f^1(x)$$

s.t.



 $\min_{\mathbf{x}\in X} f^1(\mathbf{x})$ s.t.

$\mathbf{x} \in \mathbb{R}^{n_1}$: decision vector of $\mathbf{1}^{\mathrm{st}}$ stage



$$\min_{x \in X} f^{1}(x) + \mathbb{E}[\mathcal{Q}(x, \chi)]$$
s.t.

$x \in \mathbb{R}^{n_1}$: decision vector of 1^{st} stage



$$\min_{x \in X} \quad f^1(x) + \mathbb{E}[\mathcal{Q}(x, \chi)]$$

s.t.

$$\mathcal{Q}(x,\chi) = \min_{y\in Y} f^2(x,y),$$

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$$\begin{aligned} \mathcal{Q}(x,\chi) &= \min_{y \in Y} \quad f^2(x,y), \\ \text{s.t.} \quad \mathcal{G}(x,y,\chi) \leq 0. \end{aligned}$$

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$$\min_{x \in X} f^{1}(x) + \mathbb{E}[\mathcal{Q}(x, \chi)]$$
s.t. $G^{1}(x) < 0$

$$\mathcal{Q}(x, \chi) = \min_{y \in Y} f^{2}(x, y),$$
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$$\begin{split} \min_{x \in X} & f^1(x) + \mathbb{E}[\mathcal{Q}(x,\chi)] \\ \text{s.t.} & G^1(x) < 0 \\ & \mathcal{Q}(x,\chi) = \min_{y \in Y} & f^2(x,y), \\ & \text{s.t.} & G^2(x,y,\chi) \leq 0. \end{split}$$

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Stochastic Optimization Two-Stage Programming <u>Example:</u> TSKP

Example: TSKP



Stochastic Optimization

Stochastic Optimization Two-Stage Programming Example: TSKP

Example: TSKP



Stochastic Optimization



Two-Stage Stochastic Knapsack problem





Two-Stage Stochastic Knapsack problem

First stage: items can be put in the knapsack





Two-Stage Stochastic Knapsack problem

- First stage: items can be put in the knapsack
- First stage \longleftrightarrow second stage: item weights are revealed



Two-Stage Stochastic Knapsack problem

- First stage: items can be put in the knapsack
- First stage \longleftrightarrow second stage: item weights are revealed
- Second stage: The decision can/has to be corrected



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- First stage: items can be put in the knapsack
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Second stage: Items
 ...have to be removed in case of an overweight
 ...can be added if capacity sufficient
 ...can be exchanged to increase gain.



Two-Stage Stochastic Knapsack problem

- First stage: items can be put in the knapsack
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- Second stage: Items

 ...have to be removed in case of an overweight
 ...can be added if capacity sufficient
 ...can be exchanged to increase gain.
- Correction of the decision causes penalty







- Knapsack \simeq Hotel Complex
- \blacksquare Weight capacity \simeq Total number of beds
- Items \simeq Travel groups
- Item weights \simeq Group size



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Application: Travel Agency

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- Number of beds insufficient
 - \rightarrow groups have to be relocated in other hotels



Application: Travel Agency

- Knapsack \simeq Hotel Complex
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- Items \simeq Travel groups
- Item weights \simeq Group size
- Randomness e.g., cancellations
- Agency allows overbooking
- Number of beds insufficient → groups have to be relocated in other hotels
- Vacant beds filled with last minute offers



Example: TSKP

Two-Stage Knapsack Problem

Example: TSKP



Example: TSKP

Two-Stage Knapsack Problem

$$TSKP) \max_{x \in \{0,1\}^n} \sum_{i=1}^n r_i x_i + \mathbb{E}[\mathcal{Q}(x,\chi)]$$

s.t.
$$\mathcal{Q}(x,\chi) = \max_{\mathbf{y}^+, \mathbf{y}^- \in \{0,1\}^n} \sum_{i=1}^n \overline{\mathbf{r}}_i \mathbf{y}_i^+ - \sum_{i=1}^n \mathbf{d}_i \mathbf{y}_i^-$$

 $x \in \mathbb{R}^{n}$: decision vector of 1^{st} stage $y^{+}, y^{-} \in \mathbb{R}^{n}$: decision vectors of 2^{nd} stage (recourse action) $\bar{r}_{i} < r_{i}, \ d_{i} > r_{i}$

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$$\begin{split} & x \in \mathbb{R}^n: \text{ decision vector of } 1^{\text{st}} \text{ stage} \\ & \textbf{y}^+, \textbf{y}^- \in \mathbb{R}^n: \text{ decision vectors of } 2^{\text{nd}} \text{ stage (recourse action)} \\ & \overline{\textbf{r}}_i < \textbf{r}_i, \ \textbf{d}_i > \textbf{r}_i \end{split}$$

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s.t. $y_j^+ \le 1 - x_j \quad (j = 1, \dots, n),$
 $y_i^- < x_i \qquad (j = 1, \dots, n),$

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 $\sum_{i=1}^n (x_i + y_i^+ - y_i^-) \chi_i \le c.$

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Two-Stage Programming

Example: TSKP

Linear Two-Stage Problem





 $\min_{\substack{x\geq 0\\ \text{s.t.}}} c^T x$

$x \in \mathbb{R}^{n_1}$: decision vector of 1^{st} stage



Linear Two-Stage Problem

 $\min_{x\geq 0} c^T x$ s.t. $Ax \ge b$,

$x \in \mathbb{R}^{n_1}$: decision vector of 1st stage



Linear Two-Stage Problem

$$\min_{x \ge 0} \quad c^{\mathsf{T}} x + \mathbb{E}[\mathcal{Q}(x, \chi)]$$
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Two-Stage Programming

Simple Recourse as a special case of Two-Stage problems

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Linear Simple-Recourse problem



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Linear Simple-Recourse problem

$$\min_{x\in X} \quad c^{\mathsf{T}}x + \sum_{i=1}^{m} d_i \cdot \mathbb{E}\left[[h_i(\chi) - \mathcal{T}_i(\chi)x]^+ \right]$$



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Equivalent formulation in two stages

$$\min_{\substack{x \ge 0}} c^{\mathsf{T}} x + \mathbb{E}[\mathcal{Q}(x,\chi)]$$
s.t.
$$\mathcal{Q}(x,\chi) = \min_{\substack{y \ge 0}} d^{\mathsf{T}} y$$
s.t.
$$h(\chi) - \mathsf{T}(\chi) x \le y.$$

 $x \in \mathbb{R}^n$: decision vector of 1st stage $y \in \mathbb{R}^m$: decision vectors of 2nd stage ("violation")

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Equivalent formulation in two stages

$$\min_{\substack{x \ge 0}} c^T x + \mathbb{E}[\mathcal{Q}(x,\chi)]$$
s.t.
$$\mathcal{Q}(x,\chi) = \min_{\substack{y \ge 0}} d^T y$$
s.t.
$$T(\chi)x + y \ge h(\chi).$$

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Two-Stage Programming

Simple Recourse as a special case of Two-Stage problems





Two-Stage Programming

Simple Recourse as a special case of Two-Stage problems

Linear Two-Stage Problem $\begin{array}{l} \min_{x \geq 0} \quad c^T x + \mathbb{E}[\mathcal{Q}(x, \chi)] \\ \text{s.t.} \quad Ax \geq b, \\ \mathcal{Q}(x, \chi) = \min_{y \geq 0} \quad d^T y \\ \text{s.t.} \quad T(\chi)x + W(\chi)y \geq h(\chi). \end{array}$

 $x \in \mathbb{R}^{n_1}$: decision vector of 1st stage $y \in \mathbb{R}^{n_2}$: decision vectors of 2nd stage (recourse action)

Linear Simple Recourse Problem in two stages

$$\min_{\substack{x \ge 0}} c^T x + \mathbb{E}[\mathcal{Q}(x,\chi)]$$
s.t.
$$\mathcal{Q}(x,\chi) = \min_{\substack{y \ge 0}} d^T y$$
s.t.
$$T(\chi)x + y \ge h(\chi).$$

 $x \in \mathbb{R}^n$: decision vector of 1st stage $y \in \mathbb{R}^m$: decision vectors of 2nd stage ("violation")

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— Two-Stage Programming

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Linear Simple Recourse Problem in two stages

$$\min_{\substack{x \ge 0}} \quad c^{\mathsf{T}}x + \mathbb{E}[\mathcal{Q}(x,\chi)]$$
s.t.
$$\mathcal{Q}(x,\chi) = \min_{\substack{y \ge 0}} \quad d^{\mathsf{T}}y$$
s.t.
$$\mathcal{T}(\chi)x + \mathbb{I}_n y \ge h(\chi).$$

 $x \in \mathbb{R}^n$: decision vector of 1st stage $y \in \mathbb{R}^m$: decision vectors of 2nd stage ("violation") I_n: Identity matrix of dimension *n*

Two-Stage Programming

Simple Recourse as a special case of Two-Stage problems

Simple Recourse Problem formulated as Two-Stage Problem

No first stage constraints.



Two-Stage Programming

Simple Recourse as a special case of Two-Stage problems

Simple Recourse Problem formulated as Two-Stage Problem

- No first stage constraints.
- # second stage variables = # constrains



Two-Stage Programming

Simple Recourse as a special case of Two-Stage problems

Simple Recourse Problem formulated as Two-Stage Problem

- No first stage constraints.
- # second stage variables = # constrains

$$\bullet W(\chi) = W = \mathbb{I}_n$$



- └─ Two-Stage Programming
 - Leterministic reformulation in case of discrete distributions

Outline

- 1 Randomness occurs in the constraint function
- 2 Simple Recourse problems Useful Theorems

3 Two-Stage Programming

- Example: TSKP
- Simple Recourse as a special case of Two-Stage problems
- Deterministic reformulation in case of discrete distributions
- Main theoretical result



Deterministic reformulation in case of discrete distributions

Linear Two-Stage Problem

$$\begin{split} \min_{\substack{x \geq 0}} & c^{\mathsf{T}} x + \mathbb{E}[\mathcal{Q}(x,\chi)] \\ \text{s.t.} & Ax \geq b, \\ & \mathcal{Q}(x,\chi) = \min_{\substack{y \geq 0}} & d^{\mathsf{T}} y \\ & \text{s.t.} & \mathsf{T}(\chi) x + W(\chi) y \geq h(\chi). \end{split}$$

 $x \in \mathbb{R}^{n_1}$: decision vector of 1st stage $y \in \mathbb{R}^{n_2}$: decision vectors of 2nd stage (recourse action)



Deterministic reformulation in case of discrete distributions

Linear Two-Stage Problem

$$\min_{\substack{x \ge 0}} \quad c^{\mathsf{T}}x + \mathbb{E}[\mathcal{Q}(x,\chi)]$$
s.t. $Ax \ge b$,
 $\mathcal{Q}(x,\chi) = \min_{\substack{y \ge 0}} \quad d^{\mathsf{T}}y$
s.t. $T(\chi)x + W(\chi)y \ge h(\chi)$.

 $x \in \mathbb{R}^{n_1}$: decision vector of 1st stage $y \in \mathbb{R}^{n_2}$: decision vectors of 2nd stage (recourse action) $\chi^1, \ldots, \chi^K \in \mathbb{R}^s$: scenarios $\mathbb{P}\{\chi = \chi^k\} := p^k$: probabilities



Two-Stage Programming

Deterministic reformulation in case of discrete distributions

Linear Two-Stage Problem

$$\min_{x \ge 0} \quad c^{\mathsf{T}}x + \sum_{k=1}^{\mathsf{K}} p^{k} \mathcal{Q}(x, \chi^{k})$$
s.t. $Ax \ge b$,
 $\mathcal{Q}(x, \chi) = \min_{y \ge 0} \quad d^{\mathsf{T}}y$
s.t. $T(\chi)x + W(\chi)y \ge h(\chi)$

$$\begin{split} & x \in \mathbb{R}^{n_1}: \text{ decision vector of } 1^{\text{st}} \text{ stage} \\ & y \in \mathbb{R}^{n_2}: \text{ decision vectors of } 2^{\text{nd}} \text{ stage (recourse action)} \\ & \chi^1, \dots, \chi^K \in \mathbb{R}^s: \text{ scenarios} \\ & \mathbb{P}\{\chi = \chi^k\} := p^k: \text{ probabilities} \end{split}$$

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Deterministic reformulation in case of discrete distributions

Linear Two-Stage Problem

$$\min_{\substack{x \ge 0}} c^T x + \sum_{k=1}^{K} p^k \mathcal{Q}(x, \chi^k)$$
s.t. $Ax \ge b$,
 $\mathcal{Q}(x, \chi^k) = \min_{\substack{y \ge 0}} d^T y$
s.t. $T(\chi^k)x + W(\chi^k)y \ge h(\chi^k)$.

 $x \in \mathbb{R}^{n_1}$: decision vector of 1st stage $y \in \mathbb{R}^{n_2}$: decision vectors of 2nd stage (recourse action) $\chi^1, \ldots, \chi^K \in \mathbb{R}^s$: scenarios $\mathbb{P}\{\chi = \chi^k\} := p^k$: probabilities

Deterministic reformulation in case of discrete distributions

Linear Two-Stage Problem

$$\begin{split} \min_{x \ge 0} & c^T x + \sum_{k=1}^K p^k \mathcal{Q}(x, \chi^k) \\ \text{s.t.} & Ax \ge b, \\ & \mathcal{Q}(x, \chi^k) = \min_{y \ge 0} \quad d^T y \\ & \text{s.t.} \quad T(\chi^k) x + W(\chi^k) y \ge h(\chi^k) \\ & (\forall \ k \in \{1, \dots, K\}). \end{split}$$

 $x \in \mathbb{R}^{n_1}$: decision vector of 1st stage $y \in \mathbb{R}^{n_2}$: decision vectors of 2nd stage (recourse action) $\chi^1, \ldots, \chi^K \in \mathbb{R}^s$: scenarios $\mathbb{P}\{\chi = \chi^k\} := p^k$: probabilities

Deterministic reformulation in case of discrete distributions

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s.t. $Ax \ge b$,
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 $(\forall k \in \{1, \dots, K\}).$

 $x \in \mathbb{R}^{n_1}$: decision vector of 1st stage $y \in \mathbb{R}^{n_2}$: decision vectors of 2nd stage (recourse action) $\chi^1, \ldots, \chi^K \in \mathbb{R}^s$: scenarios $\mathbb{P}\{\chi = \chi^k\} := p^k$: probabilities

Leterministic reformulation in case of discrete distributions

Deterministically Reformulated Linear Two-Stage Problem

$$\min_{x \ge 0} c^T x +$$

s.t. $Ax \ge b$

(8a)



Leterministic reformulation in case of discrete distributions

Deterministically Reformulated Linear Two-Stage Problem

$$\min_{\substack{x \ge 0 \\ y^k \ge 0}} c^T x + \sum_{k=1}^{K} p^k \left(d^T y^k \right)$$

s.t. $Ax \ge b$,

(8a)



Leterministic reformulation in case of discrete distributions

Deterministically Reformulated Linear Two-Stage Problem

$$\min_{\substack{x \ge 0 \\ y^k \ge 0}} c^T x + \sum_{k=1}^K p^k \left(d^T y^k \right)$$

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Leterministic reformulation in case of discrete distributions

Deterministically Reformulated Linear Two-Stage Problem

$$\begin{array}{ll} \min_{\substack{x \ge 0 \\ y^k \ge 0}} & c^T x + \sum_{k=1}^K p^k \left(d^T y^k \right) \\ \text{s.t.} & Ax \ge b, \\ & T(\chi^k) x + W(\chi^k) y^k \ge h(\chi^k) \quad (\forall \, k \in \{1, \dots, K\}). \end{array}$$
(8a)



Leterministic reformulation in case of discrete distributions

Deterministically Reformulated Linear Two-Stage Problem

$$\begin{array}{ll} \min_{\substack{x \ge 0\\y^k \ge 0}} & c^T x + \sum_{k=1}^K p^k \left(d^T y^k \right) \\ \text{s.t.} & Ax \ge b, \\ & T(\chi^k) x + W(\chi^k) y^k \ge h(\chi^k) \quad (\forall \, k \in \{1, \dots, K\}). \end{array}$$
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└─ Two-Stage Programming └─ Main theoretical result

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Stochastic Optimization

- Two-Stage Programming

Main theoretical result

Definition (Fixed Recourse)



Main theoretical result

Definition (Fixed Recourse)

A Linear Two Stage Programming Problem is said to have fixed recourse if the recourse matrix is deterministic, i.e. $W(\chi) = W$.



Main theoretical result

Definition (Fixed Recourse)

A Linear Two Stage Programming Problem is said to have **fixed** recourse if the recourse matrix is deterministic, i.e. $W(\chi) = W$.



Two-Stage Programming
 <u>Main theoretical result</u>

Theorem

Given: linear two-stage problem



Theorem

Given: linear two-stage problem Let:



Given: linear two-stage problem Let:

fixed recourse



└─ Two-Stage Programming └─ Main theoretical result

Theorem

Given: linear two-stage problem Let:

- fixed recourse
- $\{\lambda \in \mathbb{R}^{m_2} : d \ge \lambda W\} \neq \emptyset$



└─ Two-Stage Programming └─ Main theoretical result

Linear Two-Stage Problem

$$\min_{x \ge 0} \quad c^{\mathsf{T}} x + \mathbb{E}[\mathcal{Q}(x, \chi)]$$
s.t. $Ax \ge b,$

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 $x \in \mathbb{R}^{n_1}$: decision vector of 1st stage $y \in \mathbb{R}^{n_2}$: decision vectors of 2nd stage (recourse action) $T(\chi)$: Technology matrix $W(\chi)$: Recourse matrix



Given: linear two-stage problem Let:

- fixed recourse
- { $\lambda \in \mathbb{R}^{m_2} : d \ge \lambda W$ } $\neq \emptyset$ (\rightarrow second-stage problem primal and dual feasible)



Given: linear two-stage problem Let:

- fixed recourse
- {λ ∈ ℝ^{m₂} : d ≥ λW} ≠ Ø
 (→ second-stage problem primal and dual feasible)

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- {λ ∈ ℝ^{m₂} : d ≥ λW} ≠ Ø
 (→ second-stage problem primal and dual feasible)

Then $\mathbb{E}\left[\mathcal{Q}(\mathbf{x},\chi)\right]$ is

real-valued,



Given: linear two-stage problem Let:

- fixed recourse
- {λ ∈ ℝ^{m₂} : d ≥ λW} ≠ Ø
 (→ second-stage problem primal and dual feasible)

- real-valued,
- piecewise linear and convex in x,

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Given: linear two-stage problem Let:

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- { $\lambda \in \mathbb{R}^{m_2} : d \ge \lambda W$ } $\neq \emptyset$ (\rightarrow second-stage problem primal and dual feasible)
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Then $\mathbb{E}\left[\mathcal{Q}(\mathbf{x}, \chi)\right]$ is

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Then $\mathbb{E}\left[\mathcal{Q}(x,\chi)
ight]$ is

- real-valued,
- piecewise linear and convex in x,
- Lipschitz continuous in x,
- subdifferentiable in x,
- a polyhedral convex function.

Main theoretical result

QUESTIONS?

What about next week?

