

Stochastic Optimization

IDA PhD course 2011ht

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10. Lecture: Computational examples
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- 1 Why and when Stochastic Programming might be advantageous?
- 2 Simple example
- 3 Multiperiod Batch Plant Scheduling



Outline

- 1 Why and when Stochastic Programming might be advantageous?
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Is-Situation

- Use average/expected values
- No adaptations to current situation

Disadvantages

- Average gain \neq computed expected value
- No flexibility
- Solution not robust
- Might entail infeasible solutions



Solutions

- Online Optimization ?
- Robust Optimization
- Stochastic Programming
- ...



Typical application fields

- Capacity planning
- Energy sector
- Finance
- Forestry
- Military
- Production / Supply chain
- Scheduling
- Transportation (of humans, goods,...)
- Water management
- ...



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Example: Continuous Knapsack with uncertain capacity

$$\begin{aligned} \max_{x \in [0,1]^3} \quad & c^T x \\ \text{s.t.} \quad & a^T x \leq \chi \end{aligned}$$

- $c = (10, 15, 20)$
- $a = (5, 10, 20)$
- $p_1 = \mathbb{P}\{\chi = \chi_1\} = \mathbb{P}\{\chi = 3\} = 0.2$
- $p_2 = \mathbb{P}\{\chi = \chi_2\} = \mathbb{P}\{\chi = 12\} = 0.3$
- $p_3 = \mathbb{P}\{\chi = \chi_3\} = \mathbb{P}\{\chi = 25\} = 0.5$



Scenario dependent solutions

$$\begin{aligned} \max_{x \in [0,1]^3} \quad & c^T x(\chi) \\ \text{s.t.} \quad & a^T x(\chi) \leq \chi \end{aligned}$$

- $x(3) = (0.6; 0; 0)^T; z(3) = 6$
- $x(12) = (1; 0.7; 0)^T; z(12) = 20.5$
- $x(25) = (1; 1; 0.5)^T; z(25) = 35$
- $z = \mathbb{E}[z(\chi)] = \sum_{i=1}^3 p_i z(\chi) = 24.85$



Average value approach

$$\begin{aligned} \max_{x \in [0,1]^3} \quad & c^T x \\ \text{s.t.} \quad & a^T x \leq \mathbb{E}[\chi] \end{aligned}$$

- $x = (1; 1; 0.085)^T$
- $z = 26.7$



Worst case approach

$$\begin{aligned} \max_{x \in [0,1]^3} \quad & c^T x \\ \text{s.t.} \quad & a^T x \leq \inf_{\chi \in \Omega} \chi \end{aligned}$$

- $x = (0.6; 0; 0)^T$
- $z = 6$



Chance Constrained Approach

$$\begin{aligned} \max_{x \in [0,1]^3} \quad & c^T x \\ \text{s.t.} \quad & \mathbb{P}\{a^T x \leq \chi\} \geq \alpha \end{aligned}$$

- $\mathbb{P}\{a^T x \leq \chi\} \geq 0.8 \Leftrightarrow a^T x \leq 12$
- $x = (1; 0.7; 0)^T$
- $z = 20.5$



Two-Stage Approach

$$\begin{aligned} \max_{x \in [0,1]^3} \quad & c^T x - \sum_{i=1}^3 p_i q(\chi_i) y(\chi_i) \\ \text{s.t.} \quad & a^T x - y(\chi_i) \leq \chi_i \quad \forall i \end{aligned}$$

- $q(\chi_1) = 2; q(\chi_2) = 3; q(\chi_3) = 6$
- $x = (1; 1; 0)^T$
- $y(\chi_1) = 12$
- $y(\chi_2) = 3$
- $y(\chi_3) = 0$
- $z = 17.5$



Simple Recourse Approach

$$\max_{x \in [0,1]^3} c^T x - \mathbb{E}[q(\chi)[a^T x - \chi]^+]$$

■ $z = 17.5$



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Reference



J. Balasubramanian and I. E. Grossmann
**Approximation to Multistage Stochastic Optimization in
Multiperiod Batch Plant Scheduling under Demand
Uncertainty.** (2003)

<http://egon.cheme.cmu.edu/Papers/BalasubMultistage.pdf>



Multiproduct, multiperiod Batch production

- Different products
 - Different stages
 - Different costs (materials, stockage/holding, under-production...)
- ⇒ Hard scheduling problem

Uncertainties

- Demand
- Processing times
- Costs



Online Optimization?

- Long processing times
- Need to make predictions
- Holding/storage costs \leftrightarrow idle costs
- Difficult mathematical problem \leftrightarrow reaction times



Example Problem

- single-stage single-unit bath plant
- 2 products
- 3 processing modes (batch sizes)
- Uncertain demand
- 2 time periods (same demand distribution)



Product	Batch size (tons)	Proc. time	REV (\$/tons)	XC (\$/tons)	LC (\$/tons)
A	0-5	2	100	10	20
	5-10	4			
	10-25	6			
B	0-5	3	250	20	50
	5-10	5			
	10-25	7			

Event	Probability	Demands (tons)
1	0.25	A:10, B:0
2	0.75	A:20, B:5



Deterministic approach

- Work with expected demands over entire horizon
- Model predicts: \$5375
- Actual expected profit: \$4559



Two-Stage approach

- First-stage: production scheduling over both time periods
- Second-stage: Amount to be sold / Unsatisfied amount
- Expected profit: \$5275 (+15%)



Three-Stage approach

- First-stage: production scheduling over 1. period
- Second-stage: production scheduling over 2. period
- Third-stage: Amount to be sold / Unsatisfied amount
- Expected profit: \$5325 (+17%)



Idea

- Approximate solution
- Shrinking time horizon strategy
- Basic structure
 - Solve two-stage problem over remaining time horizon (given past demands)
 - Implement solution *only* in next time period

!

- Problem solved in advance! (\neq online optimization)
- Two-Stage problem for all nodes in scenario tree



1. Numerical Example

- single product
- 3 stages
- 3 processing modes each
- 5 time periods (different probability distributions)



Approach	Expected revenue (1000\$)	+	CPU time (sec)
Det.	656.23	0%	3
2-stage	672.04	2%	20
SH	717.32	9%	360
6-stage	722.43	10%	>50,000



2. Numerical Example

- 4 products
- 8 tasks
- 6 processing units
- 3 processing modes for each combination
- 3 time periods (different probability distributions)



Approach	Expected revenue (\$)	+	CPU time (sec)
Det.	59,573	0%	10
2-stage	62,945	6%	35
SH	75,452	27%	180
4-stage	75,851	27%	>100,000



Conclusion

- Working with average values → false estimations
- Choosing two/multi-stage model: important increase in gain
- Approximations help to decrease CPU time
- Approximations still much better than det. model



QUESTIONS?

Tack så mycket!

